

Categorified Path Calculus

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Motivation 1:

can we have string diagrams for

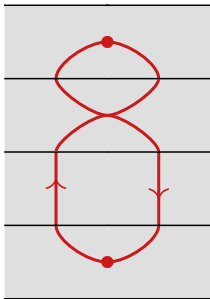
both \otimes and \oplus

?

Motivation 1: both \otimes and \oplus ?

$$\text{Tr}(A) := I \xrightarrow{\eta_A} A^* \otimes A \xrightarrow{\sigma^\otimes} A \otimes A^* \xrightarrow{\epsilon_A} I$$

$\text{Tr}(A) =$



Motivation 2:
what is an amplitude
(in quantum mechanics)

?

Path counting

Feynman Lectures on Physics, Vol. III (1965), §3-1

We wish now to put this result in terms of our new notation. First, however, we want to state our *second general principle* of quantum mechanics: When a particle can reach a given state by two possible routes, the total amplitude for the process is the *sum of the amplitudes* for the two routes considered separately. In our new notation we write that

$$\langle x | s \rangle_{\text{both holes open}} = \langle x | s \rangle_{\text{through 1}} + \langle x | s \rangle_{\text{through 2}}. \quad (3.4)$$

Parallel paths add:

$$2 + 3 = \bullet \begin{array}{c} \curvearrowright \\ \curvearrowright \\ \curvearrowright \\ \curvearrowright \\ \curvearrowright \end{array} \bullet = 5,$$

Path counting

Feynman Lectures on Physics, Vol. III (1965), §3-1

by way of hole 1. We can do that by using our *third general principle*: When a particle goes by some particular route the amplitude for that route can be written as the *product* of the *amplitude* to go part way with the *amplitude* to go the rest of the way. For the setup of Fig. 3-1 the amplitude to go from s to x by way of hole 1 is equal to the amplitude to go from s to 1, multiplied by the amplitude to go from 1 to x .

$$\langle x | s \rangle_{\text{via } 1} = \langle x | 1 \rangle \langle 1 | s \rangle. \quad (3.5)$$

Serial paths multiply:

$$2 \times 3 = \bullet \begin{array}{c} \leftarrow \text{---} \leftarrow \\ \leftarrow \text{---} \leftarrow \end{array} \bullet \begin{array}{c} \leftarrow \text{---} \leftarrow \\ \leftarrow \text{---} \leftarrow \\ \leftarrow \text{---} \leftarrow \end{array} \bullet = 6.$$

Path counting

String diagrams:

$$1 + 1 = \text{---} \bullet \text{---} \bigcirc \text{---} \bullet \text{---} = 2.$$

$$2 \times 2 = \text{---} \bullet \text{---} \bigcirc \text{---} \bullet \text{---} \bigcirc \text{---} \bullet \text{---} = \text{---} \bullet \text{---} \bigcirc \text{---} \bullet \text{---} \bigcirc \text{---} \bullet \text{---} = 4.$$

Crucial use of *distributivity*: $(a + b)c = ac + bc$.

[*Composing PROPs*, Lack, 2004]

[*Categories in Control*, Baez, Erbele, 2015]

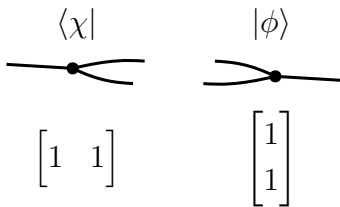
[*Interacting Hopf Algebras*, Bonchi, Sobociński, Zanasi, 2017]

[*Quantum Linear Optics via String Diagrams*, de Felice, Coecke, 2023]

Path counting

Feynman Lectures on Physics, Vol. III (1965), §8-1

One thinks of the bracket $\langle \chi | \phi \rangle$ as being divided into two pieces. The second piece $|\phi\rangle$ is often called a *ket*, and the first piece $\langle \chi |$ is called a *bra* (put together, they make a “bra-ket”—a notation proposed by Dirac); the half-symbols $|\phi\rangle$ and $\langle \chi |$ are also called *state vectors*. In any case, they are *not* numbers, and, in general, we want the results of our calculations to come out as numbers; so such “unfinished” quantities are only part-way steps in our calculations.



Path counting

$\langle \chi | \phi \rangle$



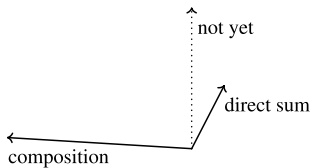
$$\begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 2 \end{bmatrix}$$

$|\phi\rangle\langle\chi|$



$$\begin{bmatrix} 1 \\ 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}$$




Path counting





$$\begin{aligned} [1] \oplus [1] &= \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} & [1] \oplus [1 \quad 1] &= \begin{bmatrix} 1 & 0 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 1 \end{bmatrix} \end{aligned}$$

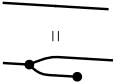
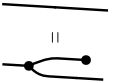
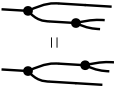
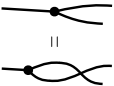
Path counting

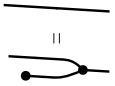
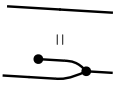
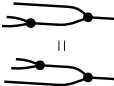
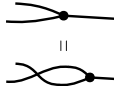
Count, unit and empty:


$$\begin{bmatrix} \\ \end{bmatrix}_{0,1}$$

$$\begin{bmatrix} \\ \end{bmatrix}_{1,0}$$

$$\begin{bmatrix} \\ \end{bmatrix}_{0,0}$$


$$\begin{bmatrix} 1 & 1 \end{bmatrix} \left(\begin{bmatrix} 1 \end{bmatrix} \oplus \begin{bmatrix} \\ \end{bmatrix}_{1,0} \right) = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \end{bmatrix}$$


Path calculus

<p>(l-unit)</p> 	<p>(r-unit)</p> 
<p>(assoc)</p> 	<p>(comm)</p> 

<p>(l-counit)</p> 	<p>(r-counit)</p> 
<p>(coassoc)</p> 	<p>(cocomm)</p> 

Path calculus

(comul-unit)		(counit-mul)	
(bimonoid)		(counit-unit)	

Punchline: equivalent to the theory of matrices with natural number entries,

$$\text{Mat}(\mathbb{N}, +, \times, 0, 1)$$

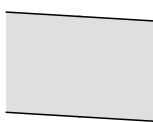
Path calculus... categorified

$$\mathbf{Mat}(\mathbb{N}, +, \times, 0, 1)$$

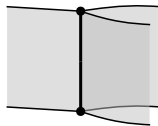
\rightsquigarrow

$$\mathbf{Mat}(\mathbf{FVec}, \oplus, \otimes, O, I)$$

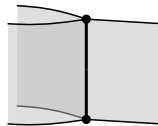
Path calculus... categorified



$[I]$

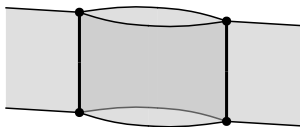


$[I \ I]$

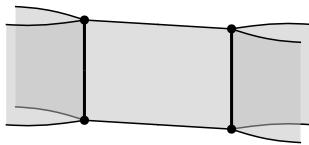


$\left[\begin{array}{c} I \\ I \end{array} \right]$

Path calculus... categorified



$$[I \ I] \begin{bmatrix} I \\ I \end{bmatrix} = [I \otimes I \oplus I \otimes I] \cong [I \oplus I]$$

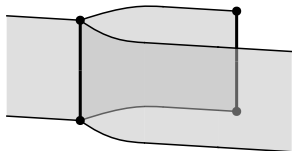


$$\begin{bmatrix} I \\ I \end{bmatrix} [I \ I] = \begin{bmatrix} I \otimes I & I \otimes I \\ I \otimes I & I \otimes I \end{bmatrix} \cong \begin{bmatrix} I & I \\ I & I \end{bmatrix}$$

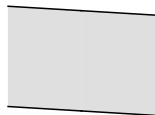
Path calculus... categorified



$[]_{1,0}$



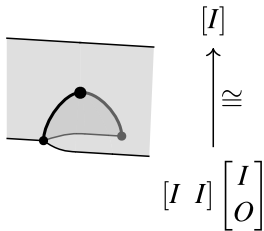
\cong



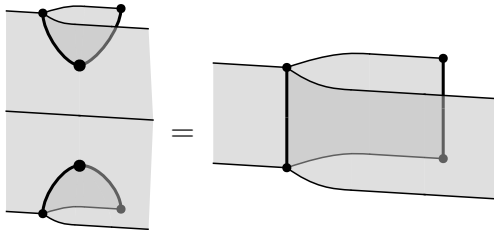
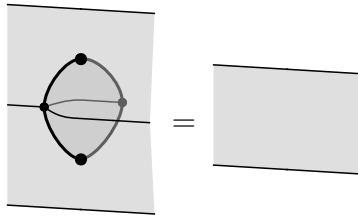
$$[I \ I] \begin{bmatrix} I \\ O \end{bmatrix} = [I \otimes I \oplus I \otimes O] \cong [I]$$

Path calculus... categorified

Coning (r-unit):



Path calculus... categorified





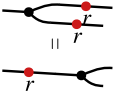
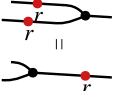
Returning to path calculus

Given any semi-ring, or rig, $(R, +, \times, 0, 1)$ we introduce the diagrammatic generators:

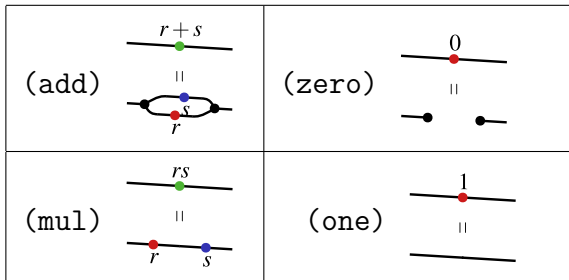


for each $r \in R$.

Returning to path calculus

<p>(unit-hom)</p> 	<p>(counit-hom)</p> 
<p>(mul-hom)</p> 	<p>(comul-hom)</p> 

Returning to path calculus



Punchline: we get a calculus for

$$\text{Mat}(R, +, \times, 0, 1)$$

Categorified path calculus

$$\text{Mat}(R, +, \times, 0, 1)$$

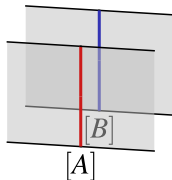
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$$\text{Mat}(\mathcal{C}, \oplus, \otimes, O, I, \alpha^\oplus, \lambda^\oplus, \rho^\oplus, \sigma^\oplus, \alpha^\oplus, \lambda^\oplus, \rho^\oplus, \dots)$$

Crucial use of *distributivity*: $(A \oplus B) \otimes C \xrightarrow{\cong} A \otimes C \oplus B \otimes C$.

[*Bimonoidal Categories, E_n -Monoidal Categories, and Algebraic K-Theory*,
Johnson, Yau, 2021]

Categorified path calculus



$$[A] \boxplus [B] := \begin{bmatrix} A & O \\ O & B \end{bmatrix}$$

Categorified path calculus

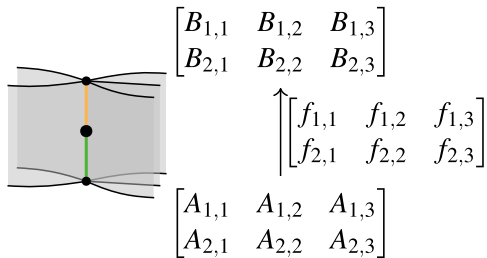
$$[f] \boxplus [g] = \begin{array}{c} \begin{array}{c} [g] \\ [f] \end{array} \\ \begin{array}{c} \bullet \\ \bullet \end{array} \end{array} \begin{array}{c} \begin{array}{c} [B \ O] \\ [O \ D] \end{array} \\ \uparrow \begin{array}{c} [f \ 1_o] \\ [1_o \ g] \end{array} \\ \begin{array}{c} [A \ O] \\ [O \ C] \end{array} \end{array}$$

Categorified path calculus

$$[f][g] = \begin{array}{ccc} [B] & [D] & [B \otimes D] \\ \uparrow [f] & \uparrow [g] & \uparrow [f \otimes g] \\ \text{---} & \text{---} & \text{---} \\ [A] & [C] & [A \otimes C] \end{array}$$

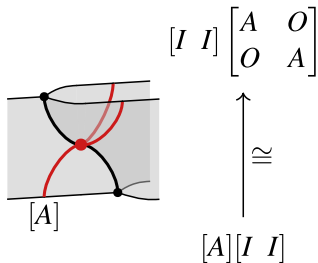
Categorified path calculus

More generally, given $f_{ij} : A_{ij} \rightarrow B_{ij}$ in \mathcal{C} , such as



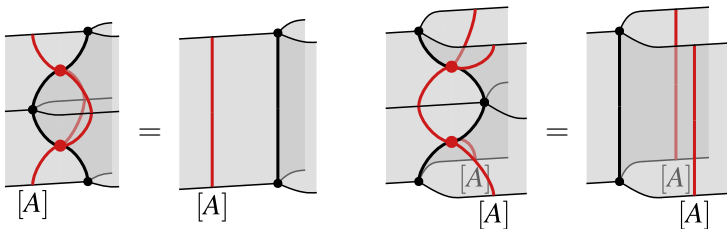
Categorified path calculus

Coning (mul-hom):



Categorified path calculus

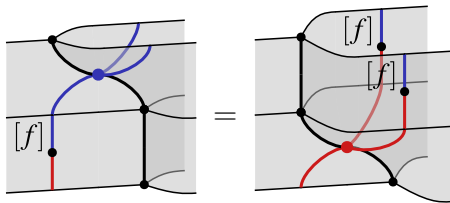
The *pulling* equations:



and *copulling*..

Categorified path calculus

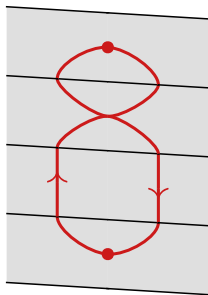
Naturality:



Backwards compatible

$$\mathrm{Tr}(A) := I \xrightarrow{\eta_A} A^* \otimes A \xrightarrow{\sigma^\otimes} A \otimes A^* \xrightarrow{\epsilon_A} I$$

$$\mathrm{Tr} \left(\begin{array}{c} \text{[Diagram: a vertical red line with a small upward-pointing arrowhead inside a gray trapezoidal box]} \\ [A] \end{array} \right) :=$$

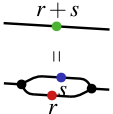
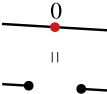
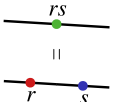
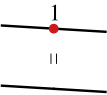


Scalars

Quantum:

$$\langle \chi | \phi \rangle \in \mathbb{C}$$

Path calculus:

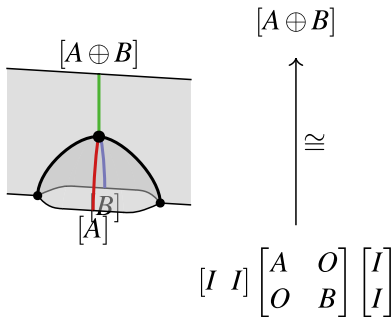
<p>(add)</p> 	<p>(zero)</p> 
<p>(mul)</p> 	<p>(one)</p> 

Categorified scalars

Hom categories:

$$\text{Hom}(1, 1) \cong \mathcal{C}$$

Categorified path calculus:

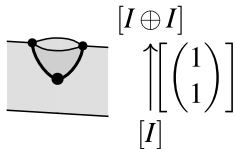
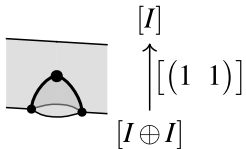


When \oplus is a Biproduct in \mathcal{C}

We have matrix morphisms in \mathcal{C} , such as

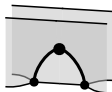
$$\begin{pmatrix} 1 & 1 \end{pmatrix} : I \oplus I \rightarrow I, \quad \begin{pmatrix} 1 \\ 1 \end{pmatrix} : I \rightarrow I \oplus I$$

these become the cap and cup:

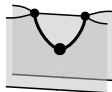


When \oplus is a Biproduct in \mathcal{C}

We also have another cap and cup



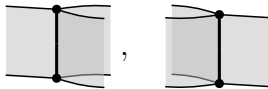
$$\begin{array}{c} \begin{bmatrix} I & O \\ O & I \end{bmatrix} \\ \uparrow \begin{bmatrix} 1 & 0_{O,I} \\ 0_{O,I} & 1 \end{bmatrix} \\ \begin{bmatrix} I & I \\ I & I \end{bmatrix} \end{array}$$



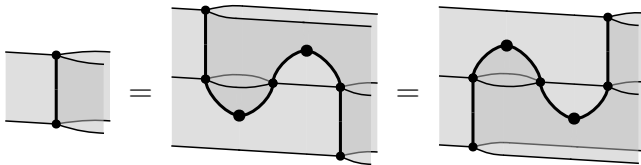
$$\begin{array}{c} \begin{bmatrix} I & I \\ I & I \end{bmatrix} \\ \uparrow \begin{bmatrix} 1 & 0_{I,O} \\ 0_{I,O} & 1 \end{bmatrix} \\ \begin{bmatrix} I & O \\ O & I \end{bmatrix} \end{array}$$

When \oplus is a Biproduct in \mathcal{C}

All these caps and cups exhibit



as ambidextrous adjoints:



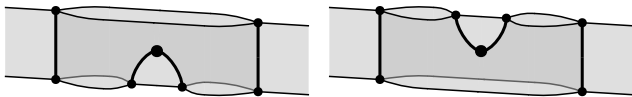
etc.

When \oplus is a Biproduct in \mathcal{C}

Ambidextrous adjoints give a Frobenius algebra on:

$$[I \oplus I] = \text{[Diagram]}$$

Multiplication, comultiplication:



... special symmetric Frobenius algebra, or *classical bit*.

[*Frobenius algebras and ambidextrous adjunctions*, Lauda, 2006]

[*Higher quantum theory*, Vicary, 2012]

2-categorical quantum mechanics

The dagger

Given $f : A \rightarrow B$ and $g : B \rightarrow D$ in \mathcal{C} :

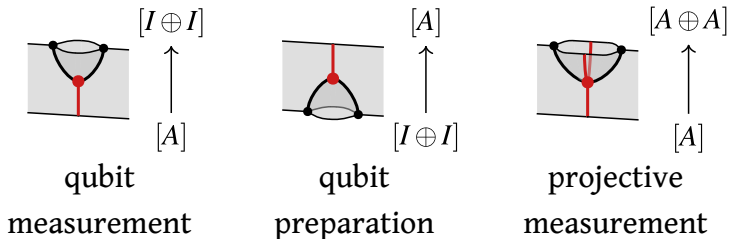
$$\left(\begin{array}{c} [C \ D] \\ \text{---} \\ [f \ g] \\ \text{---} \\ [A \ B] \end{array} \right)^\dagger := \begin{array}{c} [A \ B] \\ \text{---} \\ [f^\dagger \ g^\dagger] \\ \text{---} \\ [C \ D] \end{array} .$$

Unitary 2-cells...

[*A categorical semantics of quantum protocols*, Abramsky, Coecke, 2004]

[*Higher quantum theory*, Vicary, 2012]

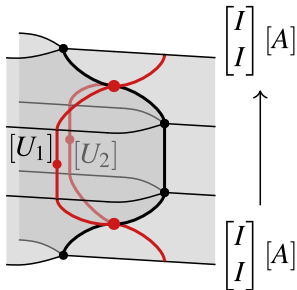
2-categorical quantum mechanics



& these are unitary.

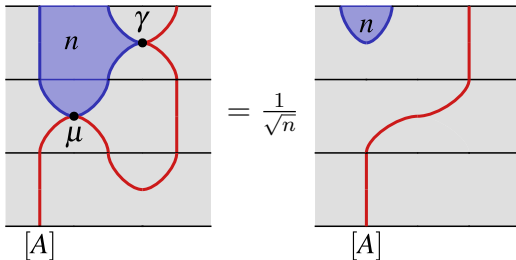
2-categorical quantum mechanics

Given unitaries $U_1, U_2 : A \rightarrow A$ in \mathcal{C} the *controlled- U* operation is the 2-cell given by

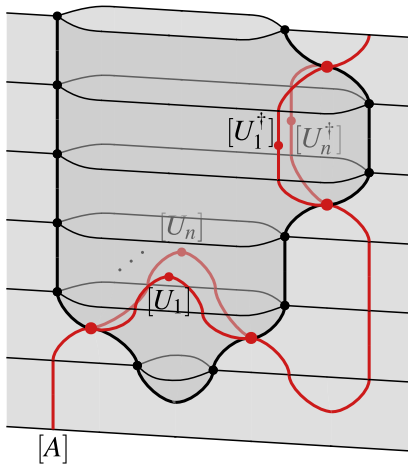


2-categorical quantum mechanics

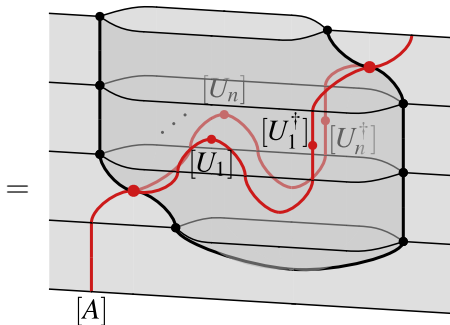
Teleportation



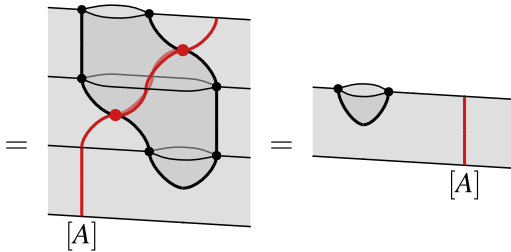
Teleportation



Teleportation



Teleportation



Conclusion

Motivation 1:

Can we have string diagrams for both \otimes and \oplus ? Yes.

Motivation 2:

What is an amplitude ? Um...

Todo:

Weak structure...

[*String diagrams for higher mathematics with wiggly.py*, Burton, 2023]

The end

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