Inference on diagrams in the category of Markov kernels

Gregoire Sergeant-Perthuis and Nils Ruet

LCQB Sorbonne Université

ACT 7 Oxford, UK 18 June, 2024

Paper arXiv:2201.11876

Sergeant-Perthuis & Ruet (LCQB)

Inference on Markov diagrams

Introduction to geometric deep learning:

- Deep learning ← curse of dimensionality
- Accounting for symmetry
 - \rightarrow Translation \rightsquigarrow CNN
- Geometry ~> discretize
 - \rightarrow Graph NN [BBCV21]
 - ightarrow Nodes share same features
 - \rightarrow Limitations: heterogeneous data
- Heterogeneity
 - \rightarrow Cellular sheaves [Cur13]
 - $\rightarrow~$ cell complex, faces \rightsquigarrow feature space, inclusions \rightsquigarrow linear maps
 - \rightarrow Functor from a poset to Vect
 - \rightarrow Sheaf Neural Networks [BGC⁺22]

First remark: limitation \rightarrow cell complexes.

- We will not talk about geometric deep learning today.
- Bayesian inference: graphical models, Markov random fields, factor graphs
- Limitation: no heterogeneity, no locality in the description of variables
- $\rightarrow\,$ Extend Bayesian inference to account for probabilistic modeling with heterogeneity and local descriptions.
 - → Independent work from PhD [SP21]

- Graphical models
- Pactor graphs
- Inference and (General) Belief Propagation
- Graphical modesl, Factor graph as contravariant functor
- **5** New!: Heterogeneous structures and probabilistic modeling
- 6 New!: Inference on diagrams in the category of Markov kernels

Definition (Undirected Graphical model)

A graphical model is the data of

- an undirected graph G = (I, A),
- a collection of variables X = (X_i ∈ E_i, i ∈ I), one per node and one variable corresponds exactly to one node

Definition (Markov properties)

Let G = (I, A) be a finite graph. Let $X = (X_i, i \in I)$ be a collection of random variables taking respectively values in the finite sets E_i . A stritcly positive probability $P_X \in \mathbb{P}(X)$ on the finite set $\Omega = \prod_{i \in I} E_i$ obeys,

(P) the pairwise Markov property relative to G, if for any pair (i, j) of non-adjacent vertices

$$X_i \perp X_j | X_{I \setminus \{i,j\}}.$$

2 (*L*) the local Markov property relative to *G*, if for any vectex $i \in V$,

$$X_i \perp X_{I \setminus (i \cup \partial i)} | X_{\partial i}$$

And we call the respective sets P(G), L(G).

Definition (Factorisation space)

Let *I* be a finite set, let $\mathscr{A} \subseteq \mathscr{P}(I)$, where $\mathscr{P}(I)$ is the set of subsets of *I*. Let $(E_i, i \in I)$ be a collection of sets, let $E_a = \prod_{i \in a} E_i$ for any $a \in \mathscr{P}(I)$; for $x \in \Omega$, we will denote x_a its projection onto E_a . The factorisation space over \mathscr{A} is defined as follows,

$$\mathsf{Fac}_{\mathscr{A}} = \{ P \in \mathbb{P}(\Omega) : \exists (f_a \in \mathbb{R}_{>0}^{E_a}, a \in \mathscr{A}), \mathsf{s. t.} \forall x \in \Omega \ P = \prod_{a \in \mathscr{A}} f_a(x_a) \}$$

$$(0.1)$$

 How to relate the Markov properties to factorizations of the underlying distribution?

Sergeant-Perthuis & Ruet (LCQB)

Definition (Cliques of a graph)

Let G = (I, A) be a graph; a clique of *G* is a subset of *G* such that every two distinct vertices are adjacent. We will note \mathscr{C} the set of its cliques.

Theorem (Hammersley-Clifford)

Let G = (I, A) be a finite graph. For all P_X strictly positive probability law on a finite set $\prod_{i \in I} E_i$,

$$P_X \in P(G) \iff P_X \in L(G) \iff P_X \in \mathsf{Fac}_{\mathscr{C}}$$
. (0.2)

• Taking the 'log': product \rightarrow sum

$$\rightarrow \prod_{a} f_{a} \rightarrow \sum_{a} H_{a}$$

 \rightarrow Relation to statistical mechanics

Inference \rightarrow Belief propagation (in few slides)

Sergeant-Perthuis & Ruet (LCQB)

- Directed graphical models: Bayesian networks
- Inference on Bayesian networks:
 - \rightarrow Define an undirected graphical model
 - \rightarrow Inference on undirected graphical model

- We want more general interaction than pairwise interaction
- Factor graphs:
 - Bipartite graphs, nodes $V = V_0 \sqcup V_1$
 - V_1 collection of $a \subseteq I$ with $a \to f_a$
 - V_0 the set of indices $i \in a$ for some $a \in V_1$
 - Edges $i \rightarrow a$ when $i \in a$

An example:

- Graphical model: X Y Z
- Factor graph: $X \to f_{X,Y} \leftarrow Y \to f_{Y,Z} \leftarrow Z$
- Factor graphs generalize graphical models
- Spaces of factorization generalize both ---- statistical mechanics.

Exemple de modèle graphique et Belief Propagation



Figure 15. An acyclic undirected graphical model.

Reference: Statistical Inference in Graphical Models, K. Gimpel, D. Rudoy



Figure 16. A message passed from node F to node E.

$$p(A) = \sum_{B} p(AB) \sum_{C} p(C \mid B) \sum_{D} p(D \mid B) m_{E \to B}(B)$$

$$= \sum_{B} p(AB) m_{E \to B}(B) \sum_{C} p(C \mid B) \sum_{D} p(D \mid B)$$

$$= \sum_{B} p(AB) m_{E \to B}(B) \sum_{C} p(C \mid B) m_{D \to B}(B)$$

$$= \sum_{B} p(AB) m_{E \to B}(B) m_{D \to B}(B) \sum_{C} p(C \mid B)$$

$$= \sum_{B} p(AB) m_{E \to B}(B) m_{D \to B}(B) m_{C \to B}(B)$$

$$= m_{B \to A}(A).$$



Figure 17. Message passes from an example run of the variable elimination algorithm to compute p(A). A node x can only send a message to a neighbor y after x has received a message from each of its neighbors (besides y). The numbers associated with each message indicate the order in which the message must be computed. If two messages have the same number, they can be computed at the same time in parallel.

Simpler case: on graphical models \rightarrow use Belief Propagation

 Belief propagation computes the marginal distributions on edges and nodes

Consider a collection of random variables (X_i, i ∈ I) and an undirected graphical model G : (I, A) that is acyclic.

$$P_{X_i,i\in I}(x_i,i\in I) = \prod_{e\in A} f_{X_e}(x_e)$$

• for each edge: two messages, $m_{X_i o X_j} \in \mathbb{P}(E_j)$

Sergeant-Perthuis & Ruet (LCQB)

Start with $m_{X_i \to X_i} = 1$ for all directed edges $(i \to j)$,

$$m_{X_i \to X_j}^{t+1}(x_j) = \sum_{x_i \in E_i} f_{\{i,j\}}(x_i, x_j) \prod_{Z \in \partial X_i \setminus X_j} m_{Z \to X_i}^t(x_i)$$
(0.3)

The stopping criteria for the algorithm is when $m_{X_i \to X_j}(x_j)^{t+1}$, which is a function over E_j , is proportional to $m_{X_i \to X_j}(x_j)^t$. Once the algorithm has finished, the marginal distributions are computed as

$$P_{X_i}(x_i) \propto \prod_{X_j \in \partial X_i} m_{X_j \to X_i}(x_i)$$
 (0.4)

Importantly, the algorithm is exact: inference is exact.

 \rightarrow For Gaussian HMM \rightsquigarrow (smoothed) Kalman filtering

Proposition (Factorization on acyclic graphs)

Let I be a finite set and let $\Omega = \prod_{i \in I} E_i$ be a product of finite sets and $X_i, i \in I$ a collection of random variables taking values respectively in E_i . Let G = (I, A) be a finite acyclic graph. $P_X \in \mathbb{P}_{>0}(E)$ factors accordingly to $\mathscr{A}(G)$, i.e., $P_X \in \operatorname{Fac}_{\mathscr{A}(G)}$ if and only if for any $\omega \in \Omega$,

$$P_X(\omega) = \frac{\prod_{e \in A} P_{X_e}(\omega_e)}{\prod_{i \in I} P_{X_i}^{d(i)-1}(\omega_i)},$$
(0.5)

where d(i) is the degree of node $i \in I$.

Sergeant-Perthuis & Ruet (LCQB)

- Bayesian inference is maximizing entropy.
- Entropy:

$$S(Q) = -\sum_{\omega \in E} Q(x) \ln Q(x)$$
(0.6)

• Bayesian inference:

 $\inf_{Q} DKL(Q \| P)$

• The same as minizing Gibbs free energy

$$\inf_{\boldsymbol{Q}\in\Theta}\mathbb{E}_{\boldsymbol{Q}}[\boldsymbol{\beta}\boldsymbol{H}]-\boldsymbol{S}(\boldsymbol{Q})$$

• But entropy:

$$\mathcal{S}(\mathcal{P}_X) = \sum_{e \in \mathcal{A}} \mathcal{S}(\mathcal{P}_{X_e}) - \sum_{i \in I} (d(i) - 1) \mathcal{S}(\mathcal{P}_{X_i})$$

- Inclusion exclusion formula c(e) = 1, c(i) = -(d(i) 1)
- Remarkably, Bayesian inference is the same as minimizing [YFW05, YFW03],

$$egin{aligned} \mathcal{F}_{\mathsf{Bethe}}(\mathcal{Q}) &= \sum_{a \in V} c(a) \mathcal{S}(\mathcal{Q}_a) - \mathbb{E}_{\mathcal{Q}_a}[\mathcal{H}_a] \end{aligned}$$

where $Q := (Q_a \in \mathbb{P}(X_a), a \in V)$ with compatibility by marginization:

 \rightarrow if *a* is an edges and *i* an edge in *a*

$$\rightarrow \pi_i^e : E_e \rightarrow E_i$$

ightarrow we ask $\pi^{e}_{i}{}_{*}(\textit{Q}_{e})=\textit{Q}_{i}$

 Belief Propagation (BP) is a discrete-time gradient descent (on Lagrange multipliers) that solves

$$\min_{Q} F_{\text{Bethe}}(Q)$$

under 'marginal' compatibility.

• Fixed points of BP correspond to critical points of *F*_{Bethe}.

 $\rightarrow\,$ I did not invent it [Pel20]... but I call it...

Definition (Graphical presheaves)

Let *I* be a finite set and $\mathscr{A} \subseteq \mathscr{P}(I)$ be a sub-poset of the powerset of *I*. Let $E_i, i \in I$ are finite sets. For $a \in \mathscr{A}$ $E_a := \prod_{i \in a} E_i$, let $F(a) := E_a$, and for $b \subseteq a$, let $F_b^a : E_a \to E_b$ be the projection map from $\prod_{i \in a} E_i$ to $\prod_{i \in b} E_i$. *F* is called a graphical presheaf from \mathscr{A} to **Mes**^{*f*}.

- Only projections
- Only products of variables, and subcollection of variables

- Graphical models
- Pactor graphs
- Inference and (General) Belief Propagation
- ④ Graphical modesl, Factor graph ... as contravariant functor
- **5** New!: Heterogeneous structures and probabilistic model
- 6 New!: Inference on diagrams in the category of Markov kernels

18/27

- Consider any map, not just projections:
 - ightarrow Measurable maps for b
 ightarrow a and even Markov kernels
- Account for possible heterogeneity, incompleteness, and incompatibility in the description of variables:
 - ightarrow Agents with different world models that communicate their beliefs
 - $\rightarrow\,$ Broader class of effective models for potential computational chemistry

19/27

- **Kern**^{*f*}: objects are finite measurable spaces, morphisms are Markov kernels (stochastic matrices).
- *F* is a contravariant functor from *A* to Kern^f; *F*^a_b : *F*(a) → *F*(b) is denoted element-wise as *F*^a_b(ω_b | ω_a), with ω_b ∈ *F*(b), ω_a ∈ *F*(a).
 - \rightarrow F encodes all the ways our data can interact.
 - $\rightarrow \mathscr{A}$ is any poset, not just a collection of subsets.
 - \rightarrow Maps are not just projections.
- $Q = (Q_a \in \mathbb{P}(F(a)), a \in \mathscr{A})$
- *F*_{Bethe}(*Q*) = ∑_{a∈𝔅} *c*(*a*) (𝔼<sub>*Q_a*[*H_a*] − *S*(*Q_a*)); *c*(*a*) = ∑_{b≥a} μ(*b*, *a*) is the generalization of the inclusion-exclusion formula associated with 𝔅.
 </sub>

For a finite poset \mathscr{A} ,

- the 'zeta-operator' of A, denoted ζ, from ⊕_{a∈A} ℝ to ⊕_{a∈A} ℝ is defined as, for any λ ∈ ⊕_{a∈A} ℝ and any a ∈ A, ζ(λ)(a) = ∑_{b≤a} λ_b
- its inverse denoted μ ; ($\mu(a, b), b \leq a$) Möbius function of \mathscr{A} .

We want to do Bayesian inference on these diagram.

- Constraint: the Q_a must be compatible under the actions of the F_b^a , i.e. $F_b^a \circ Q_a = Q_b$
- Problem: find an algorithm to 'solve' the optimization problem.
 → New message passing algorithm!

F induces several actions: on probabilities, on probabilities seen as vectors, on their dual...

- $\tilde{F}_b^a : \mathbb{P}(F(a)) \to \mathbb{P}(F(b))$ is the linear map that sends probability distributions $p \in \mathbb{P}(F(a))$ to $F_b^a \circ p$
- *˜F*^{*} is the functor obtained by dualizing the morphisms *˜F*^a_b, i.e.
 ˜F^{*,b}_a: *˜F*(*b*)^{*} → *˜F*(*a*)^{*} sends linear maps *I*_b: *˜F*(*b*) → ℝ to
 *I*_b ∘ *˜F*^a_b: *˜F*(*a*) → ℝ.

 μ can be extended to account for $\tilde{\textbf{F}},\,\tilde{\textbf{F}}^*$

 for a functor G from A to ℝ-vector spaces, we define μ_G as, for any a ∈ A and v ∈ ⊕_{a∈A} G(a), μ_G(v)(a) = ∑_{b≤a} μ(a, b)G^b_a(v_b). Recall min $F_{\text{Bethe}} = \sum_{a} F(Q_{a})$ under

- Constraint: the Q_a must be compatible under the actions of the F_b^a , i.e., $F_b^a \circ Q_a = Q_b$
 - i.e., $Q \in \lim \tilde{F}$
 - In fact, no... need to add the condition that the distribution sums to one.
 - But it's okay!

- *FE* : ∏_{a∈𝒜} P(*E_a*) → ∏_{a∈𝒜} R as *FE*(*Q*) = (E_{Q_a}[*H_a*] - *S_a*(*Q_a*), *a* ∈ 𝒜), which sends a collection of probability measures over 𝒜 to their Gibbs free energies.
- $d_Q FE \rightarrow$ differential of FE at the point Q.

Theorem

Let \mathscr{A} be a finite poset, let F be a presheaf from \mathscr{A} to **Kern**^{*f*}. Let $H_a : F(a) \to \mathbb{R}$ be a collection of (measurable) functions. The critical points of \mathscr{F} are the $Q \in \lim \tilde{F}$ such that,

$$\mu_{\tilde{F}^*} d_Q F E|_{\lim \tilde{F}} = 0 \tag{0.7}$$

Algorithm 1: Message passage algorithm for presheaves from \mathscr{A} to Kern^f

Data: Initialization: $(m_{a\to b}^0 \in \mathbb{R}^{F(b)}, b, a \in \mathscr{A} \text{ s.t. } b \leq a)$, a poset \mathscr{A} , a presheaf $F : \mathscr{A} \to \operatorname{Kern}^f$; 1 for $t \leq T$ do for $a \in \mathcal{A}, b \in \mathcal{A}$ such that $b \leq a$ do 2 $\forall \omega_a \in F(a), \quad n_{b \to a}(\omega_a) \leftarrow \prod_{\substack{c:b \leq c \\ c \neq a}} \sum_{\omega'_b \in F(b)} m_{c \to b}(\omega'_b) \cdot F^a_b(\omega'_b|\omega_a)$ 3 end 4 for $a \in \mathcal{A}, b \in \mathcal{A}$ such that $b \leq a$ do 5 $b_a = e^{-H_a} \prod_{\substack{b \in \mathscr{A}: \ b \leq a}} n_{b \to a}$ 6 $p_a = rac{b_a}{\sum_{\omega_a} b_a(\omega_a)} p_{a \to b} \leftarrow m_{a \to b} \cdot rac{ ilde{F}_b^a(p_a)}{m_a}$ 7 8 9 end 10 end

• Fix point of this message passing algorithm are critical point of *F*_{Bethe}

Sergeant-Perthuis & Ruet (LCQB)

- Michael M. Bronstein, Joan Bruna, Taco Cohen, and Petar Veličković, <u>Geometric</u> deep learning: Grids, groups, graphs, geodesics, and gauges, 2021.
- Cristian Bodnar, Francesco Di Giovanni, Benjamin Paul Chamberlain, Pietro Lio, and Michael M. Bronstein, <u>Neural sheaf diffusion: A topological perspective on</u> <u>heterophily and oversmoothing in GNNs</u>, Advances in Neural Information Processing Systems (Alice H. Oh, Alekh Agarwal, Danielle Belgrave, and Kyunghyun Cho, eds.), 2022.
- Justin Curry, <u>Sheaves, cosheaves and applications</u>, Ph.D. thesis, The University of Pennsylvania, 2013, arXiv:1303.3255.
- Olivier Peltre, <u>Message passing algorithms and homology</u>, 2020, Ph.D. thesis, Link to manuscript.
- Grégoire Sergeant-Perthuis, <u>Intersection property, interaction decomposition,</u> regionalized optimization and applications, 2021, PhD thesis, 10.13140/RG.2.2.19278.38729, Link to manuscript.

Jonathan S. Yedidia, William T. Freeman, and Yair Weiss, <u>Understanding belief</u> propagation and its generalizations, p. 239–269, Morgan Kaufmann Publishers Inc., San Francisco, CA, USA, 2003.

J.S. Yedidia, W.T. Freeman, and Y. Weiss, <u>Constructing Free-Energy</u> <u>Approximations and Generalized Belief Propagation Algorithms</u>, IEEE Transactions on Information Theory **51** (2005), no. 7, 2282–2312 (en).

Proof of Characterization of Critical Points

Understanding expression of critical points:

Zeta function ζ and Möbius functions μ for functors:

• for $u \in \bigoplus_{a \in \mathscr{A}} G(a)$, and $a \in \mathscr{A}$,

$$\zeta_G(u)(a) = \sum_{b \leq a} G_a^b(u_b)$$

$$\mu_G(u)(a) = \sum_{b \leq a} \mu(a, b) G^b_a(v_b)$$

 μ_{G} is the inverse of ζ_{G}

- 4 E b

Understanding expression of critical points:

For *F* a functor from \mathscr{A}^{op} to vector spaces, critical points *u* of 'global' regionalized loss are such that:

$$[\mu_{F^*} d_u I]|_{\lim F} = 0$$

Sergeant-Perthuis (LCQB)

・ 同 ト ・ ヨ ト ・ ヨ ト

Proof of Characterization of Critical Points

Understanding expression of critical points:

$$0
ightarrow \lim F
ightarrow igoplus_{a \in \mathscr{A}} F(a) \stackrel{\delta_F}{
ightarrow} igoplus_{a, b \in \mathscr{A} \atop a \geq b} F(b)$$

where for any $v \in \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \geq b}} F(b)$ and $a, b \in \mathscr{A}$ such that $b \leq a$, $\delta_F(v)(a,b) = F_b^a(v_a) - v_b$

This is simply stating that ker $\delta = \lim F$.

Sergeant-Perthuis (LCQB)

< 回 ト < 三 ト < 三

Proof of Characterization of Critical Points

Understanding expression of critical points:

$$0 \leftarrow (\lim F)^* \leftarrow \bigoplus_{a \in \mathscr{A}} F(a)^* \xleftarrow{\mathsf{d}_F}_{\substack{a, b \in \mathscr{A} \\ a \geq b}} F(b)^*$$

Pose d = δ^* . For any $I_{a \to b} \in \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \ge b}} F(b)^*$ and $a \in \mathscr{A}$, dm(a) = $\sum_{a > b} F_b^{a*}(m_{a \to b}) - \sum_{b > a} m_{b \to a}$

Sergeant-Perthuis (LCQB)

Rewriting condition on fix points:

 $\mu_{\textit{F}}^*\textit{d}_{\textit{u}}\textit{l} \in \mathsf{im}\,\mathsf{d}$

is the same as the fact that there is $(m_{a \rightarrow b} \in F(b)^* | a, b \in \mathscr{A}, b \leq a)$ such that,

 $d_u l = \zeta_{F^*} dm$

Sergeant-Perthuis (LCQB)

・ 同 ト ・ ヨ ト ・ ヨ

Understanding this choice of message passing algorithm:

g Lagrange multipliers *m* to $u \in \bigoplus_{a \in \mathscr{A}} F(a)$. $\delta_F(u) = 0$ defines the constraints on *u*.

 $\delta_F g \zeta_{F^*} d_F$ sends a Lagrange multiplier $m \in \bigoplus_{a,b \in \mathscr{A}} F(b)^*$ to a

constraint $c \in \bigoplus_{\substack{a,b \in \mathscr{A} \\ a \ge b}} F(b)$ defined as, for $a, b \in \mathscr{A}$ such that $b \le a$,

$$c(a,b) = \delta_F g\zeta_{F^*} d_F m(a,b) = F_b^a g_a(\zeta_{F^*} d_F m(a)) - g_b(\zeta_{F^*} d_F m(b)))$$
(0.1)

We are interested in c = 0, i.e.

$$\delta_F g \zeta_{F^*} \mathsf{d}_F m = 0$$

Sergeant-Perthuis (LCQB)

くほと くほと くほと

Understanding this choice of message passing algorithm:

Choice of algorithm on the Lagrange multipliers so that $\delta_F g \zeta_{F^*} d_F m = 0$,

$$m(t+1) - m(t) = \delta_F g \zeta_{F^*} \mathsf{d}_F m(t)$$

Any other choice would also be a good candidate!

・ 同 ト ・ ヨ ト ・ ヨ ト